Probabilistic, Modular and Scalable Inference of Typestate Specifications

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Summary

• Anek infers specifications needed to check typestate with

Plural: sound static typestate checker,
 requires pre- & post-conditions
 Specifications time-consuming

• Anek:

- Assigns and solves *probabilistic* constraints over program
 - Approximate, but works well in practice
 - Probabilities act as refinable summaries on methods
 - Scales, because it can be applied iteratively
 - Ease of design, heuristics

Plural*: An Overview

• A **sound** static typestate checker for Java programs

class File {
 // Creates an
 // open file
 File();

// File must be open
void read();

void close();

*Bierhoff, Aldrich, OOPSLA '07

Plural*: An Overview

• A **sound** static typestate checker for Java programs • Works like a type system • Assigns types to variables that vary from line to line • Requires pre- and post- types for method parameters • Type includes object state & aliasing

class File {
 // Creates an
 // open file
 File();

// File must be open
void read();

void close();

*Bierhoff, Aldrich, OOPSLA '07

Plural: An Illustration

// pre - f : Open, unique
// post - f : Closed, unique
String readAndClose(File f)

// f : Open, unique Log.log("Reading"); // f : Open, unique String s = f.read(); // f : Open, unique f.close(); // f : Closed, unique return s;

ł

2

class File {

}

// pre - this : Open, unique
// post - this : Open, unique
void read();

// pre - this : Open, unique
// post - this : Closed, unique
void close();

Plural: Other Details

• Other permissions:

Unique
Immutable
Shared
1 Writer, Many Readers

- Fields can have permissions & state
- Concurrency!
- All sound!

The Challenge

Adding Plural specifications is labor intensive

E.g., days for an API in case studies

Requires tool & methodology expertise
Requires detailed program expertise (i.e., aliasing)

Anek (अनेक): Specification Inference

- Input: Program with annotated API
- Output: Program with all annotations needed to check use of API

- Process:
 - Turn program into permission-flow graph
 - Solve probabilistic constraints
 - Inferred specification is the 'likeliest' one
- Benefits:
 - Robust to bugs
 - Easy to summarize methods, modularity
 - heuristics

Permission Graph

For all program references, create permission flow graph
Graph is almost identical to data-flow graph.
(Parameter permission implicitly returned)

Permission Graph

```
// pre - NOT GIVEN
// post - NOT GIVEN
String readAndClose(File f) {
```

```
String s = f.read();
```

...

}





Initial Probabilities

// pre - NOT GIVEN // post - NOT GIVEN String readAndClose(File f) {

String s = f.read();

....

...

class File {

```
...
// pre - this: OPEN,
// unique
void read();
}
```



	Initial Value param f pre	Initial Value read() arg pre
OPEN	?	.9999
CLOSED	?	.0001
unique	?	.9999
immutable	?	.0001
share	?	.0001
full	?	.0001
pure	?	.0001

Given known specification, how to solve for unknown?
1. Add constraints to graph
2. Off-the-shelf solver

 Given known specification, how to solve for unknown?

 Add constraints to graph
 Off-the-shelf solver

Param f (pre)		\mathbf{Z}	
	read() Arg this (pre)	2	\supset

Given known specification, how to solve for unknown?

Add constraints to graph
Off-the-shelf solver

Constraints:

Permission rules
Heuristics



Given known specification, how to solve for unknown?

Add constraints to graph
Off-the-shelf solver

Constraints:

Permission rules
Heuristics



	Final Value
OPEN	.997
CLOSED	.01
unique	.997
immutable	.01
share	.01
full	.02
pure	.03

Modularity

Solver is iterative

 Probabilities become more accurate over time

Modularity

Solver is iterative Probabilities become more accurate over time

Param f (pre)

	Intermediate Value
OPEN	.558
CLOSED	.456
unique	.677
immutable	.33
share	.25
full	.499
pure	.4

Modularity

Solver is iterative

• Probabilities become more accurate over time

• Use them as method summaries!

- 1. Solve for a method,
- 2. Store just values of signature,
- 3. Solve other methods,
- 4. Iterate

• Avoid storing entire graph

• Sliding scale between precision and time

	Intermediate Value
OPEN	.558
CLOSED	.456
unique	.677
immutable	.33
share	.25
full	.499
pure	.4

Param f (pre)

Evaluation: Training

 Evaluation Procedure:

 "Train" algorithm on small benchmarks
 Evaluate results on large case study
 Compare time & precision to manual

 • Training

- 1. Create a number of small benchmarks
- 2. Run Anek
- 3. Adjust probabilities
- 4. Iterate until results are as expected

Evaluation: Iterator API

- Given annotated Iterator API, infer specifications in PMD
 - 40kloc
 - o 170 calls to Iterator. next()
 - Initially 45 warnings
 - Case study in Bierhoff thesis

• Results

- Bierhoff:
 26 annotations in
 - **75 min**, resulted in **3 warnings**
- Anek:
 - **31 annotations** in **3 min 47 sec**, resulted in
 - 4 warnings

Evaluation: Comparisons

In paper we attempt to compare with non-probabilistic algorithm
Anek is faster because of approximative algorithms

Summary

- Plural: sound static typestate checker,
 - requires pre- & post-conditions containing aliasing permissions
 - Specifications time-consuming
- Anek infers these specifications
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 - Works well in practice
 - Probabilities act as refinable summaries on methods
 - Scales, because it can be applied iteratively
 - Ease of design, heuristics
 - We believe in probabilistic inference!
 - Source available soon!